

Solution Information Security (ME/EE)  
06-1505

Ans 1 (a) Confidentiality, Integrity, Authenticity

(b) Phishing is an internet scam where the user is convinced to give valuable information.

(c) Sniffing is a data interception technology.

Objective of Sniffing is to steal (i) Password (ii) Email text (iii) Files in transfer

(d) An attack is any action that violates security.

(e) Authentication

Authorization

→ Verifies you are who you say you are

→ Decides if you have permission to access a resource

→ Methods

a) Login form

→ Method:

b) HTTP authentication

a) Access controls for URLs

c) HTTP digest

b) Secure objects and methods

d) Custom authentication method

c) Access control lists (ACLs)

Ans 2 ⇒  $C=11$ ,  $e=7$ ,  $n=33$

$$n = p \cdot q$$

$$\therefore p=3, q=11$$

$$\phi(n) = (p-1)(q-1) = 2 \times 10 = 20$$

$$ed = 1 \pmod{\phi(n)} \Rightarrow 7 \cdot d = 1 \pmod{20}$$

$$\therefore d=3$$

Now

$$M = C^d \pmod{n}$$

$$\Rightarrow M = 11^3 \pmod{33}$$

$$\Rightarrow M = 1331 \pmod{33}$$

$$\Rightarrow M = 11 \text{ Ans}$$

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Ans 3) (a)  $2^1 \text{ mod } 11 = 2$

$2^2 \text{ mod } 11 = 4$

$2^3 \text{ mod } 11 = 8$

⋮

$2^{10} \text{ mod } 11 = 1$

∴ All the values are unique

∴ 2 is the primitive root of 11.

(b)  $g = 11, \alpha = 2$

$y_A = g, x_A = ?$

$y_A = (\alpha)^{x_A} \text{ mod } g$

$g = 2^{x_A} \text{ mod } 11$

∴  $x_A = 6$  Ans

(c)  $g = 11, \alpha = 2$

$y_B = 3$

$K = (y_B)^{x_A} \text{ mod } g$

$= (3)^6 \text{ mod } 11$

$= (729) \text{ mod } 11$

$K = 3$  Ans

Ans 4) Knapsack cryptosystem was developed by Merkle-Hellman.

The Merkle-Hellman Knapsack cryptosystem is based on NP-Complete problem.

Used for providing secure public key cryptography.

$$[ S = a_0 w_0 + a_1 w_1 + a_2 w_2 + \dots + a_{n-1} w_{n-1} ]$$

where  $a \rightarrow$  Item

$w \rightarrow$  weight,  $S =$  Sum

eg) Suppose the weights are:

85, 13, 9, 7, 47, 27, 99, 86

designated sum  $(S) = 172$

∴ Soln is  $a = (a_0, a_1, a_2, a_3, a_4, a_5, a_6, a_7) = (11001100)$

∴  $S = 85 + 13 + 47 + 27 = 172$  Ans

Ans 5 (a) IP Sec is an Authentication protocol works at NW layer

Services :- (a) Access Control

(b) Connectionless integrity

(c) Data origin Authentication

(d) Rejection of Replayed packet

(e) Confidentiality (Encryption)

(f) Limited traffic flow Confidentiality.

$$(b) C = KP \pmod{26}$$

me et me

12 4 4 19 12 4

$$a=0 \quad b=1 \quad c=2 \quad d=3 \quad e=4 \quad f=5$$

$$g=6 \quad h=7 \quad i=8 \quad j=9 \quad k=10 \quad l=11$$

$$m=12 \quad \dots \quad z=25$$

for me,

$$C = \begin{bmatrix} 9 & 4 \\ 5 & 7 \end{bmatrix} \begin{bmatrix} 12 \\ 4 \end{bmatrix} \pmod{26}$$

$$= \begin{bmatrix} 108+16 \\ 60+28 \end{bmatrix} \pmod{26} = \begin{bmatrix} 124 \\ 88 \end{bmatrix} \pmod{26} = \begin{bmatrix} 20 \\ 10 \end{bmatrix} = \begin{matrix} l \\ k \end{matrix}$$

for et

$$C = \begin{bmatrix} 9 & 4 \\ 5 & 7 \end{bmatrix} \begin{bmatrix} 4 \\ 19 \end{bmatrix} \pmod{26}$$

$$= \begin{bmatrix} 36+76 \\ 20+133 \end{bmatrix} \pmod{26} = \begin{bmatrix} 112 \\ 153 \end{bmatrix} \pmod{26} = \begin{bmatrix} 8 \\ 23 \end{bmatrix} = \begin{matrix} i \\ x \end{matrix}$$

Similarly me = (l, k)

$\therefore$  Meet me = [l k i x l k] Ans

(c) Hash Function =)

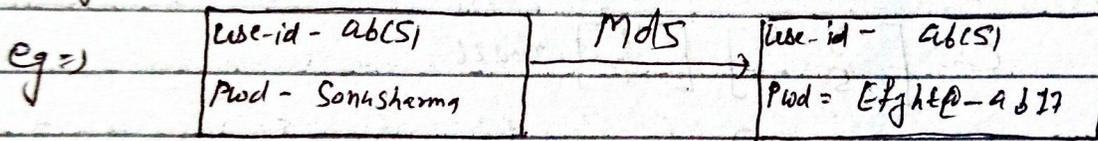
- + It's a function that takes Variable-length input string called (called pre-image) and converts it to a fixed length (generally smaller) o/p string (called hash value).
- \* A simple hash function would be a function that takes pre-image and returns a byte consisting of the XOR of all the i/c bytes.

$$h = H(m)$$

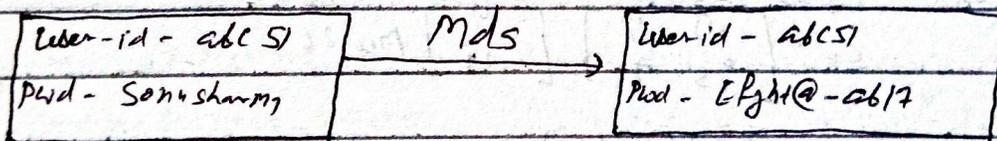
Where  $m$  is variable length msg.

$H(m)$  is the fixed length hash value.

Signup



Signin: =)



Encryption of Pwd will Save At the Server  
of Website or Service Provider Company.