

Q.1(a) CIA stands for Confidentiality, Integrity, Authenticity or Availability

Q.1(b) i) Phishing is an internet scam where users are convinced to give their valuable information.

ii) Phishing attacks can be done through mass messaging, mass mailing etc.

Q.1(c) Sniffing is a computer attack.

Sniffing is a process of monitoring and capturing all data packets passing through given network.

Q.1(d) Virus is a self duplicating program that produces its own code by attaching copies of itself into other executable codes.

(C) Authentication	Authorization
(1) Verifies user's credentials	Validates user's permissions
(2) Checks a person's details to identify him	check a user's privileges to access resources

(2) (i) RSA :- It was developed by Rivest, Shamir, Adleman.

(ii) RSA is a block cipher

RSA algorithm is as follows:-

Step 1. Key generation:-

(a) Select two prime numbers p, q where $p \neq q$ & of equal length

(b) - Calculate n

$$n = pq$$

(c) Calculate Euler totient function.

(d) Select integer 'e' such that

$$\gcd(\phi(n), e) = 1$$

$$1 < e < \phi(n)$$

(e) Calculate d

Such that

$$d = e^{-1} \pmod{\phi(n)}$$

$$\bullet \quad ed = 1 \pmod{\phi(n)}$$

(f) Now

The public key = $\{e, n\}$

Private key = $\{d, n\}$

Step 2: - Encryption

Let the plain text is capital M such that $M < n$ so cipher text

$$CT = M^e \pmod{n}$$

Step 3:-

If $CT = c$ then

Plaintext

$$PT = c \cdot \text{mod}(n)$$

We have given

$$CT = 11$$

$$e = 7$$

$$n = 33$$

Step 1 Key generation.

(a) Let $p = 3$

$$q = 11$$

$$\text{So, } n = 33$$

(b) Euler totient function

$$\phi(n) = (p-1)(q-1)$$

$$= (3-1)(11-1)$$

$$= 2 \times 10$$

$$= 20$$

(c) $ed = 1 \pmod{\phi(n)}$

$$7d = 1 \pmod{20}$$

$$d = 3$$

(d) Key generation

$$\text{Public key} = \{e, n\}$$

$$= \{7, 33\}$$

$$\text{Private key} = \{d, n\}$$

We need to find key ~~decree~~
decryption

So,

$$PT = (CT)^d \pmod{n}$$

$$= (11)^3 \pmod{33}$$

$$= 1331 \pmod{33}$$

$$= 11$$

$$\boxed{PT = 11}$$

(3) (i) Diffie-Hellman Key Exchange algorithm is a secure algorithm for public key cryptography.

(ii) The purpose of this algorithm is to enable two users to securely exchange their keys for encryption.

Diffie-Hellman algorithm:

Step 1 Global public element.

We take two no. g and a where g is prime number and a is an integer.

a is a primitive root of g .

If

$$a^1 \bmod g = x$$

$$a^2 \bmod g = y$$

$$a^3 \bmod g = z$$

$$a^n \bmod g = p$$

Where $n < g$.

Or $n = g - 1$

If we get distinct numbers like x, y, z, \dots, p then

a is primitive root of g .

Step 2 :- Public key of User A.

User A select a Random Integer

$x_A < q$ as his private

and find his public key.

$$y_A = a^{x_A} \text{ mod } q$$

Step 3 :- User B Key Generation
 $x_B < q$

$$y_B = a^{x_B} \text{ mod } q$$

Step 4:

Now A & B exchange their public key.

Step 5 :- Generation of Secret key of User A

$$K = (y_B)^{x_A} \text{ mod } q$$

Step 6 :- Generation of Secret

$$K_A = (y_A)^{x_B} \text{ mod } q$$

$$K_A = (x_A \text{ mod } q)^{x_B} \text{ mod } q$$

$$K_A = (x_B \text{ mod } q)^{x_A} \text{ mod } q$$

$$K_A = (y_A)^{x_B} \text{ mod } q$$

3.19) To show 2 is the primitive root of 11

So,

$$2^1 \text{ mod } 11 = 2$$

$$2^2 \text{ mod } 11 = 4$$

$$2^3 \text{ mod } 11 = 8$$

$$2^4 \text{ mod } 11 = 5$$

$$2^5 \text{ mod } 11 = 10$$

$$2^6 \text{ mod } 11 = 9$$

$$2^7 \text{ mod } 11 = 7$$

$$2^8 \text{ mod } 11 = 3$$

$$2^9 \text{ mod } 11 = 6$$

$$2^{10} \text{ mod } 11 = 1$$

Hence,

we get different value of mod 11 so, 2 is the primitive root of 11.

3(b) If $y_A = 9$. Find x_A

$$y_A = x^x \pmod{q}$$

$$9 = 2^{x_A} \pmod{11}$$

$$\boxed{x_A = 6} \text{ ans.}$$

3(c) $y_B = 3$ then find shared secret key.

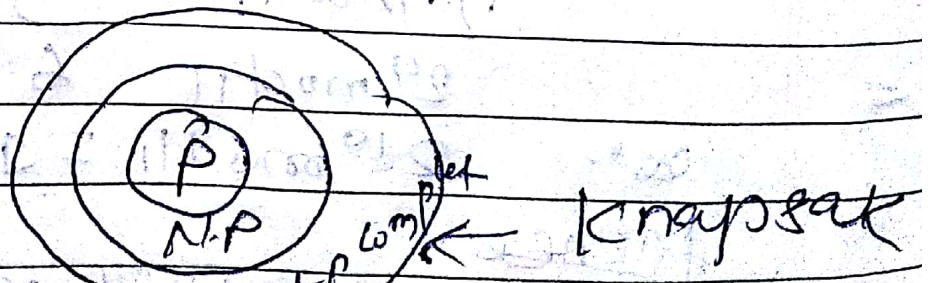
$$K = (y_B)^{x_A} \pmod{q}$$

$$= 3^6 \pmod{11}$$

$$= 729 \pmod{11}$$

$$= 3$$

(4) Knapsack crypto system is developed by Merkle Hellman use to solve $n-p$ complete problem.



→ Knapsack is used to provide secure public key system.

→ And Knapsack problem started that if we have given set of n weight $w_0, w_1, w_2, \dots, w_{n-1}$ and a sum S we have to find $a_0, a_1, a_2, \dots, a_{n-1}$ where $a \in \{0, 1\}$

$$S = a_0 w_0 + a_1 w_1 + a_2 w_2 + \dots + a_{n-1} w_{n-1}$$

where

$$S = \text{sum}$$

$a \rightarrow$ item set

$w \rightarrow$ weight

→ Generally knapsack is very difficult to used so, we are using super increasing knapsack.

A super increasing knapsack is similar to general knapsack except that where the weight are arranged from least to greatest each weight

is greater than the sum of all previous weight.

Example

If the given weight are (1, 6, 8, 15, 24) and the given sum as

(a) $S = 40$ Find the item sets.

ans) $S = 40$

$\{1, 6, 8, 15, 24\} = 2$

0	1	2	3	4
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$S > 24$

$S - 24 = 16$

$S = S - 24$

$= 40 - 24$

$= 16$

$S > 15$

$S - 15 = 1$

$S = S - 15$

$= 16 - 15$

$= 1$

$S < 8$

$$S < 6$$

$$a_1 = 0$$

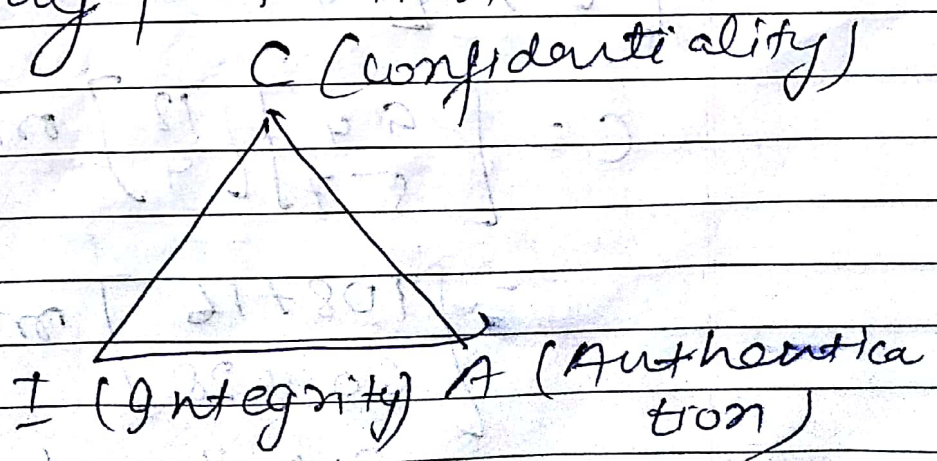
$$S > 1$$

$$a_0 = 1$$

So, item sets = $\{1, 0, 0, 1, 1\}$

5. (a) Internet protocol Security uses cryptographic security services to protect communication over internet protocol network.

IPSec supports network-level peer authentication, data origin authentication, data integrity, data confidentiality and replay protection.



5(b)

meet me

$$K = \begin{bmatrix} 9 & 4 \\ 5 & 7 \end{bmatrix}$$

$$C = K P \pmod{26}$$

ME!
ET
ME

A B C D E F G
0 1 2 3 4 5 6

H I J K L M N
7 8 9 10 11 12 13

O P Q R S T U
14 15 16 17 18 19 20

V W X Y Z
21 22 23 24 25

$$\begin{bmatrix} M \\ E \end{bmatrix} = \begin{bmatrix} 12 \\ 4 \end{bmatrix}$$

$$C = \begin{bmatrix} 9 & 4 \\ 5 & 7 \end{bmatrix} \begin{bmatrix} 12 \\ 4 \end{bmatrix} \pmod{26}$$

$$= \begin{bmatrix} 108 + 16 \\ 60 + 28 \end{bmatrix} \pmod{26}$$

$$= \begin{bmatrix} 124 \\ 88 \end{bmatrix} \pmod{26}$$

$$\begin{bmatrix} 20 \\ 10 \end{bmatrix}$$

$$\begin{bmatrix} U \\ K \end{bmatrix}$$

$$\begin{bmatrix} E \\ T \end{bmatrix} = \begin{bmatrix} 4 \\ 19 \end{bmatrix}$$

$$C = KP \pmod{26}$$

$$= \begin{bmatrix} 9 & 4 \\ 5 & 7 \end{bmatrix} \begin{bmatrix} 4 \\ 19 \end{bmatrix} \pmod{26}$$

$$= \begin{bmatrix} 36 + 76 \\ 20 + 133 \end{bmatrix} \pmod{26}$$

$$= \begin{bmatrix} 112 \\ 153 \end{bmatrix} \pmod{26}$$

$$= \begin{bmatrix} 8 \\ 1 \end{bmatrix}$$

$$= \begin{bmatrix} E \\ B \end{bmatrix}$$

$$[M] = \begin{bmatrix} 12 \\ 4 \end{bmatrix}$$

$$C = Kp \text{ mod } 26$$

$$= \begin{bmatrix} 9 & 4 \\ 5 & 7 \end{bmatrix} \begin{bmatrix} 12 \\ 4 \end{bmatrix} \text{ mod } 26$$

$$= \begin{bmatrix} 124 \\ 88 \end{bmatrix} \text{ mod } 26$$

$$= \begin{bmatrix} 20 \\ 10 \end{bmatrix}$$

$$= \begin{bmatrix} U \\ K \end{bmatrix}$$

Meet me \rightarrow U K I B U K

- (c) A hash function is a function that takes variable length input string which is called pre-image and converts it into a fixed length output called hash value.

For example:—

SHA-1 One of the most widely used cryptographic hash functions, produces a 160 bit value. In this case, the size which is called block size, is much bigger than the size of the hash value.

(6) (a) A firewall examines all the traffic related between the two networks to see if it meets certain criteria.

(b) It routes packets between the networks as well as inbound and outbound activity.

(c) It filters both inbound and outbound traffic.

Types of firewall:—

- (1) Packet filtering firewall
- (2) Network filtering firewall

Firewall

30 C

- | | |
|---|---|
| <p>① Firewall is a network security device that filters incoming and outgoing network based on pre-determined rules</p> | <p>An intrusion detection system is a device or a software application that monitors a traffic for malicious activities</p> |
| <p>② Block the traffic</p> | <p>Alarms or detection of anomaly.</p> |
| <p>③ Not analysed</p> | <p>Analysed</p> |